Supplement of “Improvement of Real-Time Multi-Core Schedulability with Forced Non-Preemption”

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DETAILS OF SECTION 2.2

The technique in [4] uses the notion of interference [27]. The interference to $\tau_k$ in an interval of $[a, b]$ (denoted by $I_k(a, b)$) represents the cumulative length of all intervals in $[a, b]$ such that a job of $\tau_k$ is ready to execute but cannot be executed due to the execution of other ready jobs. Also, the interference of a task $\tau_i$ to another task $\tau_k$ in an interval of $[a, b]$ (denoted by $I_{k-i}(a, b)$) represents the cumulative length of all intervals in $[a, b]$ such that a job of $\tau_i$ executes but a job of $\tau_k$ cannot, although it is ready for execution. Since a job of $\tau_k$ does not execute in a given time slot only when $m$ other jobs execute, the following equation holds under any global work-conserving algorithm [27]:

$$I_k(a, b) = \sum_{\tau_i \in \Phi^-(\tau_k)} \frac{I_{k-i}(a, b)}{m}. \tag{9}$$

Then, a property regarding $I_k(a, b)$ and $I_{k-i}(a, b)$ is derived, which is useful for reducing the pessimism in a schedulability analysis.

Lemma 8 (Lemma 4 in [27]): The following inequality holds for any global work-conserving algorithm:

$$I_k(a, b) \geq x \iff \sum_{\tau_i \in \Phi^-(\tau_k)} \min \left(I_{k-i}(a, b), x \right) \geq m \cdot x. \tag{10}$$

Proof: The proof can be found from [27], but is outlined here for completeness. A ready, unfinished job of $\tau_k$ does not execute in a given time slot when jobs of $m$ other tasks execute (i.e., interfere with the job of $\tau_k$) in the slot. Therefore, if we focus on the cumulative length of $x$ over all intervals in $[a, b]$ such that any job of $\tau_k$ is ready to execute but it cannot, each task’s interference with $\tau_k$ is upper-bounded by $x$, and the sum of all other tasks’ interferences with $\tau_k$ should be no less than $m \cdot x$. The opposite direction can be proved by using the definition of $I_k(a, b)$ and $I_{k-i}(a, b)$.

Using the concept of interference and Lemma 8, the technique in [4] calculates the maximum duration between the release and the completion of any job of task $\tau_k$, i.e., called the response time of $\tau_k$. We do this by computing the maximum amount of $\tau_k$’s interference with $\tau_k$ in an interval of length $l$ starting from the release of any job of $\tau_k$, which is denoted by $I_{k-l}(l)$, and formally expressed as:

$$I_{k-l}(l) \triangleq \max_{t} \left( \tau_k \in I_{k-l}(t, t + l) \right). \tag{11}$$

We define $I_{k-l}(l)$ only for $0 \leq l \leq D_k$, since we are interested in meeting the timing requirements.

If the sum of the execution time of $\tau_k$ and the maximum interference to $\tau_k$ in an interval of length $l$ starting from the release of any job of $\tau_k$ is no greater than $l$, any job of $\tau_k$ finishes its execution within $l$ time units after its release. $I_{k-l}(l)$ is used to express this, leading to the following schedulability test for fully-preemptive scheduling algorithms.

Lemma 9 (Theorem 6 in [4]): When a task set $\Phi$ is scheduled by a fully-preemptive algorithm, an upper-bound of the response time of $\tau_k \in \Phi$ is $R_k = R_k^* + 1$ such that $R_k^* \leq R_k^*$ holds in the following expression, starting from $R_k^* = C_k$:

$$R_k^{*+1} = C_k + \left[ m \cdot \sum_{\tau_i \in \Phi^-(\tau_k)} \min \left( I_{k-i}(R_k^*), R_k^* - C_k + 1 \right) \right]. \tag{12}$$

Then, if $R_k \leq D_k$ holds for all $\tau_k \in \Phi$, $\Phi$ is schedulable by the fully-preemptive algorithm.

Note that the iteration of Eq. (12) for $\tau_k$ halts if $R_k^* > D_k$, meaning that $\tau_k$ is deemed unschedulable.

Proof: The proof is given in [4], and the proof structure is the same as that of Eq. (4) in our new schedulability test of MPN scheduling algorithms to be presented in Theorem 1 in Section 4. Note that Eq. (12) is a response time analysis framework for a preemptive task since all tasks are preemptive under any fully-preemptive scheduling algorithm.

While the schedulability analysis in Lemma 9 can be applied to any preemptive (global work-conserving) scheduling algorithm, the main difficulty is to calculate $I_{k-l}(l)$ for a given scheduling algorithm. Calculating the exact $I_{k-l}(l)$ is generally intractable, and hence, the upper-bounds of $I_{k-l}(l)$ are computed. Note that when the upper-bounds are calculated, we assume that there is no deadline miss. This is because the schedulability test in [4] aims to find necessary conditions for the “first” deadline miss, as most tests do. Therefore, we will assume this for derivation of all upper-bounds of $I_{k-l}(l)$ in the rest of the paper.

The schedulability test in [4] calculates two upper-bounds of $I_{k-l}(l)$: (i) the one for any scheduling algorithm (regardless of preemptive/non-preemptive scheduling) and (ii) the other for specific preemptive scheduling algorithms (e.g., fp-EDF and fp-FP). To develop (i), we use the notion of task $\tau_i$’s slack (denoted by $S_i$) that represents the minimum interval length between the finishing time and the deadline of any job of $\tau_i$; in other words, any job of $\tau_i$ finishes its execution at least $S_i$ time units ahead of
then released and scheduled as soon as possible. In this section, we will describe how to calculate $S_i$ later in the text.

Fig. 2 shows an execution pattern of $\tau_i$'s jobs that maximizes the amount of execution of $\tau_i$'s jobs in an interval starting at $t$.

The amount of execution of a preemptive task $\tau_i$'s jobs in an interval between the release time and the deadline of $\tau_i$'s jobs is maximized as shown in Fig. 3. The number of body jobs of $\tau_i$'s jobs is calculated as (Eq. (9) in [4]):

$$I_{k-\tau_i}(l) = \min\left(I_{k-\tau_i}(l), l - C_k + 1\right)$$

so the following inequality holds under fp-EDF, for all $\tau_k \in \Phi$ and $0 \leq l \leq D_k$:

$$\sum_{\tau_i \in \Phi \setminus \{\tau_k\}} \min\left(I_{k-\tau_i}(l), l - C_k + 1\right) \leq \sum_{\tau_i \in \Phi \setminus \{\tau_k\}} \min\left(W_i(l), E_{k-\tau_i}, l - C_k + 1\right). \quad (17)$$

Now, we calculate an upper-bound of $I_{k-\tau_i}(l)$ under fp-FP. Different from fp-EDF, fp-FP maintains task-level priorities; if the priority of $\tau_i$ is higher (lower) than that of $\tau_k$, a job of $\tau_i$ can (cannot) interfere with a job of $\tau_k$, regardless of their job parameters. Therefore, $I_{k-\tau_i}(l)$ is upper-bounded by $W_i(l)$ and zero, respectively when $\tau_i$'s priority is higher and lower than $\tau_k$. Using the upper-bounds, the following inequality holds under fp-FP, for all $\tau_k \in \Phi$ and $0 \leq l \leq D_k$:

$$\sum_{\tau_i \in \Phi \setminus \{\tau_k\}} \min\left(I_{k-\tau_i}(l), l - C_k + 1\right) \leq \sum_{\tau_i \in \Phi \setminus \{\tau_k\}} \min\left(W_i(l), l - C_k + 1\right). \quad (18)$$

Then, to check the schedulability of a given task set $\Phi$ under fp-EDF or fp-FP, we use Lemma 9 with the LHS of Eq. (17) or (18) replaced with the RHS. However, one may wonder how to apply the slack tests for both fp-EDF and fp-FP. The first one is to apply Lemma 9 with the upper-bounds in Eq. (17) or (18) only once with $S_i = 0, \forall \tau_i \in \Phi$. For a tighter but higher time-complexity analysis, we briefly summarize Section 4.3 of [4]. The basic idea is to repeat the application of Lemma 9 with the upper-bounds in Eq. (17) or (18); initially, $S_i$ is set to 0 for all $\tau_i \in \Phi$, and at each iteration $S_i$ is updated by $D_i - R_i$ for all $\tau_i | R_i \leq D_i \in \Phi$. The iteration halts when there is no more update for any $S_i$ of $\tau_i \in \Phi$. We call the two schedulability tests simple and improved tests, and they correspond to Theorem 6 with Eqs. (4) and (5), and Theorem 6 with Eqs. (8) and (9) in [4], respectively.

The total time-complexity of the simple and improved tests is then $O(n^2 \cdot \max_{\tau_i \in \Phi} D_i)$ and $O(n^3 \cdot \max_{\tau_i \in \Phi} D_i^2)$, respectively [4].
Fig. 4. Generated task sets

**DETAILS OF SECTION 6.1**

There are three input parameters: (a) the task type (constrained or implicit deadlines), (b) the number of cores ($m = 2, 4, 8, 16$), and (c) individual task utilization ($C_i / T_i$) distributions (bimodal with parameter: $0.1, 0.3, 0.5, 0.7$, or $0.9$, or exponential with parameter: $0.1, 0.3, 0.5, 0.7$, or $0.9$). For each task, $T_i$ is uniformly distributed in $[1, T_{max} = 1000]$, $C_i$ is chosen based on the given bimodal or exponential parameter, and $D_i$ is uniformly distributed in $[C_i, T_i]$ for constrained deadline task systems or $D_i$ is equal to $T_i$ for implicit deadline task systems.

For each combination of (a), (b) and (c), we repeat the following procedure and generate 10,000 task sets, thus resulting in 100,000 task sets for any given $m$ and the type of task sets.

1. Initially, we generate a set of $m+1$ tasks.
2. In order to exclude unschedulable sets, we check whether the generated task set can pass a necessary feasibility condition [28].
3. If it fails to pass the feasibility test, we discard the generated task set and return to Step 1. Otherwise, we include this set for evaluation. Then, this task set serves as a basis for the next new set; we create a new set by adding a new task into an already created and tested set, and return to Step 2.

Then, Fig. 4 shows the total number of generated task sets with different task set utilization (i.e., $U_{sys} = \sum_{i \in \Phi} C_i / T_i$) in $[U_{sys} - 0.01 \cdot m, U_{sys} + 0.01 \cdot m]$. Among the generated task sets, Figs. 5 and 6 to be presented plot the number of task sets proven schedulable by each schedulability analysis, with task set utilization in $[U_{sys} - 0.01 \cdot m, U_{sys} + 0.01 \cdot m]$. In all figures including

6. For a given bimodal parameter $p$, a value for $C_i / T_i$ is uniformly chosen in $[0.05]$, with probability $p$, and in $[0.5, 1]$ with probability $1-p$.
7. For a given exponential parameter $1/\lambda$, a value for $C_i / T_i$ is chosen according to an exponential distribution whose probability density function is $\lambda \cdot \exp(-\lambda \cdot x)$.

Figs. 4, 5 and 6, we selectively choose $m$ values since the task set generation or simulation results with different values of $m$ have similar behaviors.

**DETAILS OF SECTION 6.2**

Fig. 5 illustrates the average schedulability results under EDF with different preemption policies, which were discussed in Section 6.2.

To show the effect of the number of tasks on the schedulability improvement by mpn-EDF, we focus on constrained deadline task sets with $m = 4$ and $1.5 \leq U_{sys} \leq 2.0$, and investigate two statistics over different numbers of tasks as shown in Table 4: (i) the ratio of the difference between the number of schedulable task sets under mpn-EDF and fp-EDF, to the number of schedulable task sets under fp-EDF (denoted by $\%$ of non-preemptive tasks), and (ii) the average number of non-preemptive tasks of each task set under mpn-EDF when the task set is schedulable with mpn-EDF but not with fp-EDF (denoted by # of non-preemptive tasks).

As shown in the table, mpn-EDF enables each task set to be schedulable by making the small number of tasks non-preemptive (i.e., no larger than 1.13 in any case in the table). This is because disallowing preemption of each task makes a negative impact on other tasks’ schedulability; if the preemptions of more tasks are
disallowed, the other tasks that were originally schedu-
alable with fp-EDF are more likely to be unschedulable. 
Therefore, mpn-EDF is less effective on task sets each of
whose number of unschedulable tasks with fp-EDF is
larger.

On the other hand, the interference-based schedulabil-
ity framework in Section 2.2 is known to have greater
pessimism with more tasks due to over-estimation of
each task’s interference from carry-in jobs [26]. Therefore,
the more tasks in each task set, the more room for
improvement.

Due to these two conflicting effects, the improvement
by mpn-EDF increases with the number of tasks, and
peaks with 9 and 10 tasks (in which the second factor
dominates). After that, the improvement decreases (in
which the first factor dominates).

**DETAILS OF SECTION 6.3**

Similar to the EDF case, we also show the number of task
sets schedulable by the following schedulability tests of
FP with different preemption policies:

- The only existing np-FP tests [6, 8] (denoted by np-
FP[GG]), i.e., schedulable by at least one of the two
tests;
- Our schedulability test for np-FP, i.e., Theorem 1
with the upper-bounds in Lemma 4 when $Y_i = 0$
for all $\tau_i \in \Phi$ (denoted by np-FP);
- The existing fp-FP test [4], which is equivalent to
our schedulability test for fp-FP, i.e., Theorem 1 with
the upper-bounds in Lemma 4 when $Y_i = 1$ for all
$\tau_i \in \Phi$ (denoted by fp-FP);
- np-FP and fp-FP (denoted by np-FP+fp-FP), i.e.,
schedulable by at least one of np-FP and fp-FP;
- Our schedulability test of mpn-FP with disal-
lowance preemption by Algorithm 2 when the initial
preemption requirement is $Y_i = 1$ for all $\tau_i \in \Phi$
(denoted by mpn-FP).

As in Section 6.2, we present our schedulability tests
with the best performance, i.e., applying slack reclama-
tion. For fixed-priority of tasks, we deploy DM (Deadline
Monotonic) [29], in which the priority of a task $\tau_k$
is higher than that of another task $\tau_i$ if $D_k < D_i$.

Fig. 6 plots the number of schedulable task sets by
each schedulable test with different $m$ and the type of
task sets. If we focus on the schedulability under np-FP,
np-FP finds 11.3%, 17.3%, 27.2% and 54.5% additional
schedulable constrained deadline task sets (likewise
26.7%, 34.6%, 39.5% and 46.3% additional schedulable

![Fig. 5. Schedulability results under EDF with different preemption policies](image-url)
The number of schedulable task sets

Task set utilization

- mpn-FP
- np-FP + fp-FP
- fp-FP
- np-FP
- np-FP[GG]

(a) $m = 4$, constrained deadline task sets

(b) $m = 16$, constrained deadline task sets

(c) $m = 4$, implicit deadline task sets

(d) $m = 16$, implicit deadline task sets

Fig. 6. Schedulability results under FP with different preemption policies

implicit deadline task sets), which are not schedulable by np-FP[GG].

The schedulability trend between mpn-FP and np-FP+fp-FP is similar to that between mpn-EDF and np-EDF+fp-EDF. That is, the number of additional schedulable task sets covered by mpn-FP is significant, and the number is increasing as the number of cores becomes larger. The numeric values of task sets which are schedulable by mpn-FP but not by np-FP+fp-FP are 5.3%, 16.0%, 30.9% and 47.6% with implicit task sets (8.7%, 20.5%, 32.6% and 44.3% with constrained task sets) for $m = 2, 4, 8$ and 16, respectively.